Adaptive Learning in Continuous Games: **Optimal Regret Bounds and Convergence to Nash Equilibrium**

Yu-Guan Hsieh¹, Kimon Antonakopoulos^{1,2}, Panayotis Mertikopoulos^{1,2,3}

Online Learning in Continuous Games

- At each round t = 1, 2, ..., each player $i \in \mathcal{N} := \{1, ..., n\}$
- Plays an action $x_t^i \in \mathcal{X}^i$
- Suffers loss $\ell^i(\mathbf{x}_t)$ and receives as feedback $g_t^i :=$
- Joint action of all players $\mathbf{x} = (x^i)_{i \in \mathcal{N}} = (x^i, \mathbf{x}^{-i})$
- $\ell^i(\cdot, \mathbf{x}^{-i})$ is convex and $\nabla_i \ell^i(\mathbf{x}_t)$ is Lipschitz continuou
- Nash equilibrium \mathbf{x}_{\star} : $\forall i \in \mathcal{N}, \forall x^i \in \mathcal{X}^i, \ell^i(x^i_{\star}, \mathbf{x}^{-i}_{\star}) \leq \mathcal{X}^i$
- Individual regret of agent *i*:

$$\operatorname{Reg}_{T}^{i}(\mathcal{P}^{i}) = \max_{p^{i} \in \mathcal{P}^{i}} \sum_{t=1}^{I} \left(\underbrace{\ell^{i}(x_{t}^{i}, \mathbf{x}_{t}^{i}, \mathbf{$$

$$\underbrace{\ell \left(\underbrace{\ell^{i}(x_{t}^{i}, \mathbf{x}_{t}^{-i}) - \ell^{i}(p^{i}, \mathbf{x}_{t}^{-i})}_{\text{cost of not playing } p^{i} \text{ in round } i} \right)}_{i}$$

Limitations of Existing Methods

Need for adaptive stepsize

Two players run optimistic gradient with learning rates (η_t) for $\ell^{1}(\mathbf{x}) = -\ell^{2}(\mathbf{x}) = x^{1}x^{2}; \quad \mathcal{X}^{1} = \mathcal{X}^{2} = [-4, 8]; \quad x_{\star} = (0, 0)$



Need for Dual Averaging update

Feedback should enter the algorithm with equal weight

Assume linear loss and simplex-constrained action set.

- $\mathcal{X} = \{(w_1, w_2) \in \mathbb{R}^2_+, w_1 + w_2 = 1\}$
- Feedback sequence:

$$\underbrace{-e_1, \ldots, -e_1}_{\lceil T/3 \rceil}, \underbrace{-e_2, \ldots, -e_2}_{\lfloor 2T/3 \rfloor}$$

• Algorithm: Adaptive (Optimistic) Multiplicative Weight Update



TL;DR

We introduce for learning in continuous games a family of algorithms that are Adaptive, No-regret, Consistent, and Convergent.

Optimistic Dual Averaging



Regularizer h^i : 1-strongly convex and Mirror map: $Q^{i}(y) = \arg \max_{x \in \mathcal{X}^{i}} \langle y, x \rangle$ – Bregman divergence: $D^{i}(p, x) = h^{i}(p)$

Energy inequality and adaptive learning rate

Fenchel coupling. $F^i(p, y) = h^i(p) + (h^i)^*(y) - \langle y, p \rangle$ Let $\lambda_t^i = 1/\eta_t^i$, $\psi_t^i(p^i) = F^i(p^i, Y_t^i)$, $\varphi^i(p^i) = h^i(p^i) - \min h^i$. Then, for any $p^i \in \mathcal{X}^i$, we have

 $\lambda_{t+1}^{i} \psi_{t+1}^{i}(p^{i}) \leq \lambda_{t}^{i} \psi_{t}^{i}(p^{i}) - \langle g_{t}^{i}, X_{t+\frac{1}{2}}^{i} - p^{i} \rangle +$ convergence measure + $\langle g_{t}^{i} - g_{t-1}^{i}, X_{t+\frac{1}{2}}^{i} - X_{t+1}^{i} \rangle - \lambda$

gradeint variation

Rearranging, we get

 $\sum_{t=1}^{T} \langle g_{t}^{i}, X_{t+\frac{1}{2}}^{i} - p^{i} \rangle \leq \lambda_{T+1}^{i} \varphi^{i}(p^{i}) + \sum_{t=1}^{T} \left[\frac{\|g_{t}^{i} - y_{t}\|}{\|g_{t}^{i} - y_{t}\|} \right]$

Take the adaptive learning rate

$$\eta_t^i = \frac{1}{\sqrt{\tau^i + \sum_{s=1}^{t-1} ||g_t^i|}}$$

• $\tau^i > 0$ can be chosen freely by the player • η_t^i is thus computed solely based on local information available to each player

$$\nabla_i \ell^i(\mathbf{x}_t)$$

US
$$\ell^i(x^i, \mathbf{x}_{\star}^{-i})$$

$$\underbrace{t}_{t}^{-i}$$
)).



(¹UGA, Inria ²CNRS ³Criteo Al Lab)

$$= -\eta_t^i \sum_{s=1}^{t-1} g_t^s$$

= $Q(Y_t^i) = \operatorname*{arg\,min}_{x \in \mathcal{X}^i} \sum_{s=1}^{t-1} \langle g_s^i, x \rangle + \frac{h^i(x)}{\eta_t^i}$
 $+ \frac{1}{2} = \operatorname*{arg\,min}_{x \in \mathcal{X}^i} \langle g_{t-1}^i, x \rangle + \frac{D^i(x, X_t^i)}{\eta_t^i}$

$$C^{1}$$

$$- h^{i}(x)$$

$$- h^{i}(x) - \langle \nabla h^{i}(x), p - x \rangle$$

$$(\lambda_{t+1}^i - \lambda_t^i)\varphi^i(p^i)$$

$$\lambda_t^i D^i(X_{t+1}^i, X_{t+\frac{1}{2}}^i) - \lambda_t^i D^i(X_{t+\frac{1}{2}}^i, X_t^i)$$

distance between successive iterates

$$\frac{g_{t-1}^{i}\|_{(i),*}^{2}}{\lambda_{t}^{i}} - \sum_{t=2}^{T} \frac{\lambda_{t-1}^{i}}{8} \|X_{t+\frac{1}{2}}^{i} - X_{t-\frac{1}{2}}^{i}\|_{(i)}^{2}$$

(Adapt)

quence incurs $\mathcal{O}(\sqrt{T})$ regret.

verges to best response $\arg \min_{x^i \in \mathcal{X}^i} \ell^i(x^i, \mathbf{x}_{\infty}^{-i})$.

a. The game is strictly variationally stable.

$$\langle \mathbf{V}(\mathbf{x}), \mathbf{x} - \mathbf{x}_{\star} \rangle \coloneqq \sum_{i=1}^{N} \langle \nabla \mathbf{v}_{i} \rangle$$

a strict inequality whenever $\mathbf{x} \notin \mathcal{X}_{\star}$.

Proof sketch for convergence to NE

- monotonicity of the iterates with respect to any NE.
- 1. Show that λ_t^i converges to a finite constant when $t \to +\infty$. 2. Under a suitable divergence metric, establish the quasi-Fejér
- 3. Derive that $\|\mathbf{X}_{t+\frac{1}{2}} \mathbf{X}_t\| \to 0$ and $\|\mathbf{X}_t \mathbf{X}_{t-\frac{1}{2}}\| \to 0$ as $t \to +\infty$. 4. Prove that every cluster point of the sequence of play is a NE
- and conclude.





Theoretical Guarantees

- No-regret: Playing the algorithm against any bounded feedback se-
- **Consistent:** If \mathcal{X}^i is compact and the action profile \mathbf{x}_t^{-i} of all other players converges to some limit profile \mathbf{x}_{∞}^{-i} , the algorithm itself *i* con-
- Convergent: When employed by all players, the induced sequence of play converges to a Nash equilbrium if either
- b. The game is variationally stable and h^i is (sub)differentiable on all \mathcal{X}^i .
- Variational Stability. A convex game is variationally stable if the set \mathcal{X}_{\star} of Nash equilibria of the game is nonempty and
 - $\nabla_i \ell^i(\mathbf{x}), x^i x^i_{\star} \geq 0$ for all $\mathbf{x} \in \mathcal{X}, \mathbf{x}_{\star} \in \mathcal{X}_{\star}$.
- The game is strictly variationally stable if the inequality holds as
- Examples: Convex-concave zero-sum Zero-sum polymatrix